CSC 2421H : Graphs, Matrices, and Opt	imization	Lecture 6 :	October 22, 2018
Electrical Networks and their	Applications to	Random	Walks
Lecturer: Sushant Sachdeva		Scribe:	Calum MacRury

## **1** Electrical Networks

Let us suppose that G = (V, E) is an undirected graph, for which |V| = n and |E| = m. It will be convenient to assume in this section that its edges are of the form  $\{u, v\}$  for  $u, v \in V$ , highlighting the fact that they are undirected. We shall later replace each edge  $e = \{u, v\} \in E$ , with one of (u, v) or (v, u) to indicate the direction we wish to assign to e. In the former case, we say that u is directed to v, and in the latter case, we say that v is directed to u.

If we additionally specify a pair of distinct vertices  $s, t \in V$ , then we may consider an *electrical* current on G. Formally, this is a vector  $\mathbf{f} \in \mathbb{R}^{E}$ . Intuitively, the current assigned to an edge should also have a direction associated to it. In order to specify these directions, we first assign an arbitrary orientation to the edges of G. If  $\{u, v\} \in E$ , then we can orient (direct) it from u to v by replacing it with the directed edge (u, v), and by writing  $u \to v$ . We then adopt the following convention for determining the direction of the  $|f_e|$  units of current for the (directed) edge  $e = (u, v) \in \tilde{E}$ :

- If  $f_{u,v} > 0$ , then we say that  $|f_{u,v}|$  units of current move from u to v.
- If  $f_{u,v} < 0$ , then we that  $|f_{u,v}|$  units of current move from v to u.
- If  $f_{u,v} = 0$ , then 0 units of current move in either direction.

Using this convention, the vector  $\mathbf{f} \in \mathbb{R}^E$  is sufficient to describe the directions of the electrical current on G.

Now that we have formalized our notation for representing electrical currents, we consider a number of *conservation conditions* known as Kirchoff's Laws. Intuitively, these laws say that for each *internal node* (i.e  $v \neq s, t$ ), the total amount of current entering v is equal to the total amount of current leaving it. That is, formally we have

$$\sum_{u:v \to u} f_{u,v} - \sum_{w:w \to v} f_{w,v} = 0,$$

for each  $v \in V$ ,  $v \neq s, t$ .

Of course, if we consider s to be the *source* of the circuit, and t to be its *sink*, then we can generalize these constraints to include s and t as well. To do so, let us suppose that  $\gamma \ge 0$  is the total external current entering the source s. In this case, we have that

$$\sum_{u:v \to u} f_{u,v} - \sum_{w:w \to v} f_{w,v} = \begin{cases} \gamma & \text{if } v = s \\ -\gamma & \text{if } v = t \\ 0 & \text{otherwise} \end{cases}$$

For our purposes, we shall normalize and assume that the amount of current entering s is 1; that is,  $\gamma = 1$ . In this case, there exists a unique matrix  $\mathbf{B} \in \mathbb{R}^{n \times m}$  encoding the constraints we impose on **f**. We have that **f** satisfies the conservation laws, if and only if

$$\mathbf{Bf} = \mathbb{1}_s - \mathbb{1}_t$$

It is easy to check that for  $(u, v) \in E$ , if  $\mathbf{b}_{u,v} := \mathbb{1}_u - \mathbb{1}_v$ , then

$$\mathbf{B} = [\mathbf{b}_{e_1} \dots \mathbf{b}_{e_m}],$$

provided  $e_1, \ldots, e_m$  are the (directed) edges of G.

Let us now additionally associate a *resistance*  $r_e$  to each edge  $e \in E$ . If we consider the vector  $\mathbf{r} := (r_e)_{e \in E}$ , then we can define the diagonal matrix  $\mathbf{R} \in \mathbb{R}^{m \times m}$ , by setting  $\mathbf{Diag}(\mathbf{R}) = \mathbf{r}$ . We can use  $\mathbf{R}$  to define the *energy* of the current  $\mathbf{f}$ . This is defined to be,

$$\sum_{e \in E} \frac{1}{2} r_e f_e^2 = \frac{1}{2} \mathbf{f}^T \mathbf{R} \mathbf{f}$$

With this definition, we can consider a natural optimization problem, in which **b** is taken to be a fixed vector within  $\mathbb{R}^V$ , and **B** is the matrix defined above:

minimize 
$$\frac{1}{2} \mathbf{f}^T \mathbf{R} \mathbf{f}$$
  
subject to  $\mathbf{B} \mathbf{f} = \mathbf{b}$  (1)  
 $\mathbf{f} \in \mathbb{R}^E$ 

We say that a current  $\mathbf{f}$  is an *electrical current*, provided it is an optimum solution to OP 1.

**Proposition 1.1.** If **f** is an optimum solution to OP 1 then **Rf** is orthogonal to  $\Delta \in \mathbb{R}^m$ , provided  $\mathbf{B}\Delta = 0$ .

*Proof.* Assume  $\mathbf{B}\Delta = 0$ . Given  $\epsilon \in \mathbb{R}$ , we know that  $\mathbf{f} + \epsilon \Delta$  is a feasible solution to OP 1, as

$$\mathbf{B}(\mathbf{f} + \epsilon \Delta) = \mathbf{B}\mathbf{f} + \epsilon \mathbf{B}\Delta = \mathbf{B}\mathbf{f} = b.$$

On the other hand, we know that  $\mathbf{f}$  is an optimum solution so,

$$\frac{1}{2}(\mathbf{f} + \epsilon \Delta)^T \mathbf{R}(\mathbf{f} + \epsilon \Delta) \ge \frac{1}{2} \mathbf{f}^T \mathbf{R} \mathbf{f}.$$

After simplification, it follows that

$$\epsilon(\Delta^T \mathbf{R} \mathbf{f} + \epsilon \frac{1}{2} \Delta^T \mathbf{R} \Delta) \ge 0, \tag{2}$$

for all  $\epsilon$ . Notice also that the entries of **R** are nonnegative, so  $\Delta^T \mathbf{R} \Delta \ge 0$ . We consider two cases:  $\epsilon > 0$  and  $\epsilon < 0$ . First, consider  $\epsilon > 0$ . Then,

$$\Delta^T \mathbf{R} \mathbf{f} + \epsilon \frac{1}{2} \Delta^T \mathbf{R} \Delta \ge 0.$$
(3)

Now assume  $\Delta^T \mathbf{R} \mathbf{f} < 0$ . Then, let  $\epsilon_0 = -\frac{2\Delta^T \mathbf{R} \mathbf{f}}{\Delta^T \mathbf{R} \Delta}$ . Note that  $\epsilon_0 > 0$  since  $\Delta^T \mathbf{R} \mathbf{f} < 0$ . This means that

$$\epsilon < \epsilon_0 \Rightarrow \Delta^T \mathbf{R} \mathbf{f} + \epsilon \frac{1}{2} \Delta^T \mathbf{R} \Delta < 0,$$

which is a contradiction. So, the assumption must be false and it must be true that  $\Delta^T \mathbf{R} \mathbf{f} \ge 0$ .

Now consider the case where  $\epsilon < 0$ . Then,

$$\Delta^T \mathbf{R} \mathbf{f} + \epsilon \frac{1}{2} \Delta^T \mathbf{R} \Delta \le 0.$$
(4)

Now assume  $\Delta^T \mathbf{R} \mathbf{f} > 0$ . Then, let  $\epsilon_0 = -\frac{2\Delta^T \mathbf{R} \mathbf{f}}{\Delta^T \mathbf{R} \Delta}$ . Note that  $\epsilon_0 < 0$  since  $\Delta^T \mathbf{R} \mathbf{f} > 0$ . This means that

$$\epsilon > \epsilon_0 \Rightarrow \Delta^T \mathbf{R} \mathbf{f} + \epsilon \frac{1}{2} \Delta^T \mathbf{R} \Delta > 0,$$

which is a contradiction. So, the assumption must be false and it must be true that  $\Delta^T \mathbf{R} \mathbf{f} \leq 0$ .

From the above cases, the only valid possibility is that  $\Delta^T \mathbf{R} \mathbf{f} = 0$ , proving that  $\mathbf{R} \mathbf{f}$  is indeed orthogonal to  $\Delta$ .

Suppose that we now consider the subspace  $C := \{\Delta \in \mathbb{R}^E : \mathbf{B}\Delta = 0\}$ . Clearly, C is the kernel of the matrix **B** (which is typically denoted by ker(**B**)). Moreover, as a corollary to the above proposition, we get the following result:

**Corollary 1.2.** If **f** is an optimal solution to Equation 2, then **Rf** is orthogonal to the kernel of **B**; that is,  $\mathbf{f} \in C^{\perp}$  in the above notation.

We shall now consider some properties of the subspace C. Before doing so, let us observe a proposition regarding the matrix **B**.

**Proposition 1.3.** We have that  $rank(\mathbf{B}) = n - 1$  if and only if G is connected (where we consider G as an undirected graph).

*Proof.* It is a standard result from linear algebra that both the row and column spaces of  $\mathbf{B}$  have the same dimension. This integer is defined precisely as the value of rank( $\mathbf{B}$ ).

As a consequence of this result, we know that since  $\mathbf{B}^T$  is the transpose of  $\mathbf{B}$ , each matrix must have the same rank. It is therefore sufficient to show that the result holds for the matrix  $\mathbf{B}^T$ .

Observe that for any vector  $\mathbf{x} \in \mathbb{R}^V$ , we have that

$$(\mathbf{B}^T)_e = x_u - x_v,$$

for each directed edge e = (u, v) of G. In particular, this shows that  $\mathbf{B}^T \mathbb{1}_V = 0$ , and so the kernel of  $\mathbf{B}^T$  has dimension greater or equal to 1.

One can additionally show that if G is connected, then provided  $\mathbf{B}^T \mathbf{x} = 0$ , we must have that  $\mathbf{x} \in \text{span}\{\mathbb{1}_V\}$ . In other words, the dimension of the kernel of  $\mathbf{B}^T$  is exactly 1. In this case, the rank theorem for matrices implies that  $\text{rank}(\mathbf{B}^T) = n - 1$ .

Generalizing the above results, we can prove that if  $k(G) \ge 1$  is the number of components of G, then we have that rank $(\mathbf{B}^T) = n - k(G)$  (here the components are formed considering G as an undirected graph). This completes both directions of the proposition.

**Lemma 1.4.** If G is connected, then dim(C) = m - (n - 1).

*Proof.* This lemma is an immediate corollary of the previous result, combined with the rank theorem for matrices.

**Lemma 1.5.** If  $Im(\mathbf{B}^T) := {\mathbf{B}^T \mathbf{x} : \mathbf{x} \in \mathbb{R}^V}$ , namely the image of the matrix  $\mathbf{B}^T$ , then  $Im(\mathbf{B}^T) = C^{\perp} = ker(B)^{\perp}$ .

Before we prove this lemma, we remark that this statement is a property of all matrices - it is not specific to  $\mathbf{B}$ .

*Proof.* Observe that  $\text{Im}(\mathbf{B}^T)$  is spanned by the column vectors of  $\mathbf{B}^T$ . As a result, a vector  $\mathbf{x} \in \mathbb{R}^E$  is in  $\text{Im}(\mathbf{B})^{\perp}$  if and only if it is orthogonal to all the column vectors of  $\mathbf{B}^T$ . On the other hand, we know that  $\mathbf{x}$  satisfies this property if and only if it is in the kernel of  $\mathbf{B}$ . We may therefore conclude that,

$$\operatorname{Im}(\mathbf{B}^T)^{\perp} = C$$

Taking the orthogonal complement of both sides of the above equation, we may conclude that

$$\operatorname{Im}(\mathbf{B}^T) = C^{\perp},$$

thus completing the proof.

As a result of the preceding lemmas, we may conclude the following theorem:

**Theorem 1.6.** If **f** is an optimum solution to OP 1, then there exists some  $\mathbf{x} \in \mathbb{R}^E$  such that  $\mathbf{f} = \mathbf{R}^{-1}\mathbf{B}^T\mathbf{x}$ .

*Proof.* Observe that since **f** is an optimum solution, we know that **Rf** is orthogonal to C. On the other hand, we know that  $\text{Im}(B^T) = C^{\perp}$ , so there exists some  $\mathbf{x} \in \mathbb{R}^V$  for which  $\mathbf{Rf} = \mathbf{B}^T \mathbf{x}$ . As the matrix **R** is invertible, the result thus holds.

We know that if  $\mathbf{f}$  is an optimum solution to OP 1, then we have that

$$\mathbf{b} = \mathbf{B}\mathbf{f} = \mathbf{B}\mathbf{R}^{-1}\mathbf{B}^T\mathbf{x},$$

as  $\mathbf{B}\mathbf{f} = \mathbf{b}$  and  $\mathbf{f} = \mathbf{R}^{-1}\mathbf{B}^T\mathbf{x}$ . We can interpret the vector  $\mathbf{x}$  as specifying a *voltage* of  $x_v$  units for each vertex  $v \in V$ . By applying  $\mathbf{R}^{-1}\mathbf{B}^T$  to this vector  $\mathbf{x}$ , we can then recover the electrical current  $\mathbf{f}$  on the edges G, given that the resistances are defined by  $\mathbf{R}$ .

In the next section, we shall characterize exactly what these voltage vectors look like. Before we discuss how this can be done, we first observe that the matrix  $\mathbf{BR}^{-1}\mathbf{B}^T$  corresponds to the Laplacian of a specific weighted graph we now describe.

**Lemma 1.7.** The matrix  $\mathbf{BR}^{-1}\mathbf{B}^T$  is the Laplacian of G = (V, E, w), where  $w_e := \frac{1}{r_e}$  for each  $e \in E$ .

*Proof.* Observe that if R = I, the identity matrix, then

$$\mathbf{B}\mathbf{B}^{T} = \sum_{e=(u,v)\in E} (\mathbb{1}_{u} - \mathbb{1}_{v})(\mathbb{1}_{u} - \mathbb{1}_{v})^{T}$$
$$= \sum_{e=(u,v)\in E} L_{u,v},$$

where  $L_{u,v}$  is the Laplacian of the edge  $e = (u, v) \in E$ . If the matrix R is not the identity, then a similar computation shows that

$$\mathbf{B}\mathbf{R}^{-1}\mathbf{B}^{T} = \sum_{e=(u,v)\in E} \frac{1}{r_e} (\mathbb{1}_u - \mathbb{1}_v) (\mathbb{1}_u - \mathbb{1}_v)^{T}$$
$$= \sum_{e=(u,v)\in E} \frac{1}{r_e} L_{u,v},$$

where the matrix  $\frac{1}{r_e}L_{u,v}$  is the Laplacian of the edge e = (u, v) with weight  $\frac{1}{r_e}$ .

## 1.1 Moore-Pensoose Pseudo-Inverse

Suppose that we are given a symmetric  $n \times n$  matrix **A** together with an *n*-vector **b**, and we wish to solve the equation  $\mathbf{A}\mathbf{x} = \mathbf{b}$  for solve variable  $\mathbf{x} \in \mathbb{R}^n$ . In the case in which **A** is invertible, the vector  $\mathbf{x} = \mathbf{A}^{-1}\mathbf{b}$  is the unique solution to this equation. When the matrix **A** is not invertible, then we can define the *Moore-Pensoose Pseudo-Inverse* of **A**. If  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of **A**, then assume that  $\psi_1, \ldots, \psi_n$  are orthonormal eigenvectors of **A**. Using the spectral decomposition of **A**, we know that

$$\mathbf{A} = \sum_{i=1}^{n} \lambda_i \psi_i \psi_i^T.$$

We then define

$$\mathbf{A}^+ := \sum_{i=1:\lambda_i \neq 0}^n \frac{1}{\lambda_i} \psi_i \psi_i^T,$$

as the *pseudo-inverse* of **A**.

If we consider the specific case when we are given an undirected graph G = (V, E), then we can consider the *pseudo-inverse*  $\mathbf{L}^+$  of its Laplacian  $\mathbf{L}$ . In this case, assuming that  $\lambda_1 \leq \ldots \leq \lambda_n$  are the eigenvalues of  $\mathbf{L}$ , then we have that

$$\mathbf{L}(\mathbf{L}^{+}\mathbf{b}) = \sum_{i=1:\lambda_{i}>0}^{n} \psi_{i}\psi_{i}^{T}\mathbf{b},$$
(5)

where  $\psi_1, \ldots, \psi_n$  are orthonormal eigenvectors of  $\lambda_1, \ldots, \lambda_n$ . Observe that if we define the matrix  $\Pi := \sum_{i=1:\lambda_i>0}^n \psi_i \psi_i^T$ , then  $\Pi \in \mathbb{R}^{V \times V}$ . Moreover,  $\Pi$  is an orthogonal projection onto

the subspace spanned by  $\beta := \{\psi_i : \lambda_i > 0 \text{ and } 1 \leq i \leq n\}$ . That is,  $\Pi^2 = \Pi$ ,  $\Pi^T = \Pi$  and  $\operatorname{Im}(\Pi) = \operatorname{span}(\beta)$ .

For our purposes, we are particularly interested in the case when G is connected. If this is true, then we know that the kernel of **L** is spanned by  $\mathbb{1}_V$ , and so  $\lambda_i > 0$  for i = 2, ..., n. If we also assume that **b** is orthogonal to  $\mathbb{1}_V$ , then  $\mathbf{b} \in \text{span}(\beta)$ , and so  $\Pi(\mathbf{b}) = \mathbf{b}$ .

Under these assumptions, observe that Equation 5 simplifies to

$$\mathbf{L}(\mathbf{L}^+\mathbf{b}) = \mathbf{b}$$

We may therefore conclude that  $\mathbf{x} = \mathbf{L}^+ \mathbf{b}$  is a solution to the equation " $\mathbf{L}\mathbf{x} = \mathbf{b}$ ".

**Lemma 1.8.** If G = (V, E) is a connected graph, and  $\mathbf{b} \in \mathbb{R}^V$  is orthogonal to  $\mathbb{1}_V$ , then

$$\{\mathbf{x} \in \mathbb{R}^V : \mathbf{L}\mathbf{x} = \mathbf{b}\} = \{\mathbf{L}^+\mathbf{b} + \alpha \mathbb{1}_V : \alpha \in \mathbb{R}\},\$$

**Remark 1.9.** If we orient the edges of G and consider currents on the edges of G, then the condition on **b** has a natural interpretation: We can think of the assumption  $\mathbf{b}^T \mathbb{1}_V = 0$  as enforcing the constraint that the net current into the circuit must be 0. For example, in the case of a single source-sink pair (s,t), the vector  $\mathbf{b} := \mathbb{1}_s - \mathbb{1}_t$  has exactly one unit of current entering s and one unit of current leaving t. Clearly, the orthogonality condition is satisfied in this case.

*Proof.* Observe that if  $\alpha \in \mathbb{R}$ , then we have that

$$\mathbf{L}(\mathbf{L}^+\mathbf{b} + \alpha \mathbb{1}_V) = \mathbf{b},$$

as  $\mathbb{1}_V$  is in the kernel of **L**. We may therefore conclude that

$$\{\mathbf{L}^+\mathbf{b} + \alpha \mathbb{1}_V : \alpha \in \mathbb{R}\} \subseteq \{\mathbf{x} \in \mathbb{R}^V : \mathbf{L}\mathbf{x} = \mathbf{b}\}.$$

To see the other inclusion, assume that  $\mathbf{x} \in \mathbb{R}^V$  is such that  $\mathbf{L}\mathbf{x} = \mathbf{b}$ . In this case, we have that  $\mathbf{L}\mathbf{x} = \mathbf{b} = \mathbf{L}(\mathbf{L}^+\mathbf{b})$ . Thus,

$$\mathbf{L}(\mathbf{L}^+\mathbf{b} - \mathbf{x}) = 0,$$

and so  $\mathbf{L}^+\mathbf{b} - \mathbf{x} \in \ker(\mathbf{L})$ . But G was assumed to be connected, so the kernel of  $\mathbf{L}$  is spanned by  $\mathbb{1}_V$ . It follows that there exists some  $\alpha_0 \in \mathbb{R}$  such that  $\mathbf{L}^+\mathbf{b} - \mathbf{x} = \alpha_0 \mathbb{1}_V$ . Thus,  $\mathbf{x} = \mathbf{L}^+\mathbf{b} - \alpha_0 \mathbb{1}_V$ , and so  $\mathbf{x} \in {\mathbf{L}^+\mathbf{b} + \alpha \mathbb{1}_V : \alpha \in \mathbb{R}}$ . This implies the other direction of the inclusion, and so the statement holds.

We conclude this section by remarking that if we orient the graph G as in the previous section, and specify a resistance matrix  $\mathbf{R}$  on its edges, then we can consider the special matrix defined by  $\mathbf{L} = \mathbf{B}\mathbf{R}^{-1}\mathbf{B}^T$ , where the matrix  $\mathbf{B}$  is derived from the orientation on G. As we saw previously,  $\mathbf{L}$ is in fact a Laplacian matrix. If we set  $\mathbf{b} := \mathbb{1}_s - \mathbb{1}_t$  for some source-sink pair  $(s, t) \in V$ , then we can interpret this vector as passing 1 unit of current into s, and 1 unit out of t. In particular, we know that  $\mathbb{1}^T \mathbf{b} = 0$ . Thus, if we consider solutions to the equation " $\mathbf{L}\mathbf{x} = \mathbf{b}$ ", we know that

$$\{\mathbf{x} \in \mathbb{R}^V : \mathbf{L}\mathbf{x} = \mathbf{b}\} = \{\mathbf{L}^+\mathbf{b} + \alpha \mathbb{1}_V : \alpha \in \mathbb{R}\},\$$

by Lemma 1.8. Observe that if  $\mathbf{x} := \mathbf{L}^+ \mathbf{b} + \alpha \mathbb{1}_V$  for some  $\alpha \in \mathbb{R}$ , then  $\mathbf{x}$  specifies voltages on the vertices of G. Moreover, by applying  $\mathbf{R}^{-1}\mathbf{B}^T$  to  $\mathbf{x}$ , we can recover the current

$$\mathbf{f} = \mathbf{R}^{-1} \mathbf{B}^T (\mathbf{L}^+ \mathbf{b} + \alpha \mathbb{1}_V).$$

Of course,  $\mathbf{B}^T \mathbb{1}_V = 0$ , so this means that  $\mathbf{f} = \mathbf{R}^{-1} \mathbf{B}^T (\mathbf{L}^+ \mathbf{b})$ . In particular, the current  $\mathbf{f}$  must be an optimum solution to OP 1 by Theorem 1.6. In other words, any voltage solution of the above form induces an electrical current through G; that is, a current whose energy is minimum.

## 1.2 Random Walks

Let us consider a weighted undirected graph  $G = (V, E, \mathbf{w})$ , together with a pair of distinct nodes  $s, t \in V$ . If we assume that G is connected, and start a (non-lazy) random walk at node s, then we can define h(s,t) to be the expected number of steps for the walk to reach t for the first time. One can show that this value is finite, no matter which nodes are chosen.

If we fix the node t, then we can define the vector  $\mathbf{h}_t \in \mathbb{R}^V$ , where

$$\mathbf{h}_t(s) := h(s, t),$$

for each  $s \in V(G)$ . We first observe that for  $x \neq t$ ,

$$\mathbf{h}_t(x) = 1 + \sum_{y:\{x,y\}\in E} \frac{w_{x,y}}{\deg(x)} \mathbf{h}_t(y),$$

and  $\mathbf{h}_t(t) = 0$ . As a consequence, we know that for each  $x \neq t$ ,

$$\deg(x)\mathbf{h}_t(x) = \deg(x) + \sum_{y:\{x,y\}\in E} w_{x,y}\mathbf{h}_t(y).$$

In vector notation, provided  $x \neq t$ , we can express this as

$$(\mathbf{Dh}_t)(x) = (\mathbf{D}\mathbb{1}_V)(x) + (\mathbf{Ah}_t)(x),$$

provided  $\mathbf{A}$  is the adjacency matrix of G, and  $\mathbf{D}$  is its degree matrix. Observe then that

$$(\mathbf{Lh}_t)(x) = (\mathbf{D}\mathbb{1}_V)(x),$$

for all  $x \neq t$ , where **L** is the Laplacian of G. On the other hand, we know that

$$0 = \mathbb{1}_V^T \mathbf{L} \mathbf{h}_t = \sum_{x \neq t} (\mathbf{L} \mathbf{h}_t)(x) + (\mathbf{L} \mathbf{h}_t)(t),$$

as  $\mathbb{1}_V$  is in the kernel of L, and is thus orthogonal to  $\mathbf{Lh}_t$  (check this). It follows that

$$(\mathbf{L}\mathbf{h}_t)(t) = -\sum_{x \neq t} \deg(x) = \deg(t) - \mathbb{1}^T \mathbf{D} \mathbb{1}.$$

If we define  $\mathbf{b} := \mathbf{D} \mathbb{1}_V - (\mathbb{1}^T \mathbf{D} \mathbb{1}) \mathbb{1}_t$ , then this implies that

$$\mathbf{L}\mathbf{h}_t = \mathbf{b}$$

We may therefore use Lemma 1.8 to conclude that

$$\mathbf{h}_t = \mathbf{L}^+ (\mathbf{D}\mathbb{1} - (\mathbb{1}^T \mathbf{D}\mathbb{1})\mathbb{1}_t) + \alpha \mathbb{1},$$

for some  $\alpha \in \mathbb{R}$ , as **b** is orthogonal to 1. Observing that  $\mathbf{h}_t(t) = 0$ , we may conclude that

$$\mathbb{1}_t^T \mathbf{h}_t = \mathbb{1}_t^T (\mathbf{L}^+ \mathbf{b} + \alpha \mathbb{1}) = 0.$$

Thus,  $\alpha = -\mathbb{1}_t^T \mathbf{L}^+ \mathbf{b}$ . If we fix some  $s \in V$ , then this implies that

$$h(s,t) = \mathbb{1}_s^T \mathbf{h}_t$$
  
=  $\mathbb{1}_s^T \mathbf{L}^+ \mathbf{b} + \alpha \mathbb{1}_s^T \mathbb{1}$   
=  $\mathbb{1}_s^T \mathbf{L}^+ \mathbf{b} - \mathbb{1}_t^T \mathbf{L}^+ \mathbf{b}$   
=  $(\mathbb{1}_s - \mathbb{1}_t)^T \mathbf{L}^+ (\mathbf{D}\mathbb{1} - (\mathbb{1}^T \mathbf{D}\mathbb{1})\mathbb{1}_t),$ 

giving us a convenient expression for the hitting time vector  $\mathbf{h}_t$ .

## 1.3 Commute Time

In addition to hitting times on undirected weighted graphs, we can also define commute times. If  $G = (V, E, \mathbf{w})$  is connected, then provided  $s, t \in V$ , we can define the commute time from s to t as

$$C(s,t) := h(s,t) + h(t,s).$$

By our results from the previous section, we know that

$$h(s,t) = (\mathbb{1}_s - \mathbb{1}_t)^T \mathbf{L}^+ (\mathbf{D}\mathbb{1} - (\mathbb{1}^T \mathbf{D}\mathbb{1})\mathbb{1}_t)$$

Thus, after simplification

$$C(s,t) = [(\mathbb{1}_s - \mathbb{1}_t)^T \mathbf{L}^+ (\mathbb{1}_s - \mathbb{1}_t)] (\mathbb{1}^T \mathbf{D} \mathbb{1}).$$

We remark that the vector  $\mathbf{L}^+(\mathbb{1}_s - \mathbb{1}_t)$  can be interpreted as specifying voltages on the vertices of G. If 1 unit of current enters s and 1 unit leaves t, then an electrical current can be derived from this vector (see the end of the previous section for details).